

Solutions - Midterm Exam

(February 18th @ 7:30 pm)

Presentation and clarity are very important! Show your procedure!

PROBLEM 1 (17 PTS)

- Compute the result of the following operations. The operands are signed fixed-point numbers. The result must be a signed fixed-point number. For the division, use $x = 4$ fractional bits.

10.011 + 0.11111	1.010101 - 01.0101
1.10101 × 10.111	10.0110 ÷ 0.0111

$$\begin{array}{r}
 \begin{array}{cccccccc}
 \text{0} & \text{0} & \text{1} & \text{1} & \text{1} & \text{0} & \text{0} & \text{0} \\
 \text{1} & 0.0 & 1 & 1 & 0 & 0 & + \\
 0 & 0.1 & 1 & 1 & 1 & 1 & \\
 \hline
 1 & 1.0 & 1 & 0 & 1 & 1 &
 \end{array}
 &
 \begin{array}{r}
 \begin{array}{cccccccc}
 \text{1} & 1.0 & 1 & 0 & 1 & 0 & 1 & - \\
 0 & 1.0 & 1 & 0 & 1 & 0 & 0 & \\
 \hline
 1 & 1.0 & 1 & 0 & 1 & 0 & 1 & + \\
 1 & 0.1 & 0 & 1 & 1 & 0 & 0 & \\
 \hline
 1 & 0.0 & 0 & 0 & 0 & 0 & 1 &
 \end{array}
 \end{array}
 \end{array}$$

$$\begin{array}{r}
 \begin{array}{r}
 1.10101 \times \\
 10.111
 \end{array}
 \rightarrow
 \begin{array}{r}
 0.01011 \times \\
 01.001
 \end{array}
 \rightarrow
 \begin{array}{r}
 1011 \times \\
 1001 \\
 \hline
 1011 \\
 0000 \\
 0000 \\
 0000 \\
 \hline
 1011 \\
 01100011 \\
 \hline
 0.01100011
 \end{array}$$

- ✓ $\frac{10.0110}{0.0111}$: To unsigned and then alignment, $a = 4$: $\frac{01.1010}{0.0111} = \frac{01.1010}{0.0111} \equiv \frac{11010}{111}$

$$\begin{array}{r}
 000111011 \\
 111 \overline{) 110100000} \\
 \underline{111} \\
 1100 \\
 \underline{111} \\
 1010 \\
 \underline{111} \\
 1100 \\
 \underline{111} \\
 1010 \\
 \underline{111} \\
 11
 \end{array}$$

Append $x = 4$ zeros: $\frac{110100000}{111}$

Integer Division:

$$\begin{aligned}
 Q &= 111011, R = 11 \\
 \rightarrow Qf &= 11.1011 (x = 4)
 \end{aligned}$$

$$\text{Final result (2C): } \frac{10.0110}{0.0111} = 2C(011.1011) = 100.0101$$

PROBLEM 2 (8 PTS)

- Represent these numbers in Fixed Point Arithmetic (signed numbers). Use the FX format [12 4].

✓ -17.125

✓ 13.75

✓ -17.125:

$$17.125 = 010001.001 \rightarrow -17.125 = 101110.111 = 0xEE.E$$

✓ 13.75:

$$13.75 = 01101.11 = 0xD.C$$

PROBLEM 3 (40 PTS)

- Perform the following 32-bit floating point operations. For fixed-point division, use 4 fractional bits. Truncate the result when required. Show your work: how you got the significand and the biased exponent bits of the result. Provide the 32-bit result.

✓ C3FA8000 - C1E00000	✓ D0D80000 + D0FA0000	✓ 80C00000 × FAD00000	✓ 7B380000 ÷ C8A00000
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- ✓ $X = \text{C3FA8000} - \text{C1E00000}$:

C3FA8000: 1100 0011 1111 1010 1000 0000 0000 0000

$$e + \text{bias} = 10000111 = 135 \rightarrow e = 135 - 127 = 8$$

$$\text{Significand} = 1.11110101$$

$$\text{C3FA8000} = -1.11110101 \times 2^8$$

C1E00000: 1100 0001 1110 0000 0000 0000 0000 0000

$$e + \text{bias} = 10000011 = 131 \rightarrow e = 131 - 127 = 4$$

$$\text{Significand} = 1.11$$

$$\text{C1E00000} = -1.11 \times 2^4$$

$$X = -1.11110101 \times 2^8 + 1.11 \times 2^4 = -1.11110101 \times 2^8 + \frac{1.11}{2^4} \times 2^8$$

$$X = -(1.11110101 - 0.000111) \times 2^8$$

$$\text{We perform unsigned subtraction: } X = -1.11011001 \times 2^8$$

$$X = -1.11011001 \times 2^8, e + \text{bias} = 8 + 127 = 135 = 10000111$$

$$X = 1100 0011 1110 1100 1000 0000 0000 0000 = \text{C3EC8000}$$

$$\begin{array}{r} \begin{array}{cccccccc} b_9 & b_8 & b_7 & b_6 & b_5 & b_4 & b_3 & b_2 & b_1 & b_0 \end{array} \\ \begin{array}{cccccccccccc} 1 & 1 & 1 & 1 & 1 & 0 & 1 & 0 & 1 & - \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 & 0 & \\ \hline 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 & 1 & \end{array} \end{array}$$

- ✓ $X = \text{D0D80000} + \text{D0FA0000}$:

D0D80000: 1101 0000 1101 1000 0000 0000 0000 0000

$$e + \text{bias} = 10100001 = 161 \rightarrow e = 161 - 127 = 34$$

$$\text{Significand} = 1.1011$$

$$\text{D0D80000} = -1.1011 \times 2^{34}$$

D0FA0000: 1101 0000 1111 1010 0000 0000 0000 0000

$$e + \text{bias} = 10100001 = 161 \rightarrow e = 161 - 127 = 34$$

$$\text{Significand} = 1.1111010$$

$$\text{D0FA0000} = -1.1111010 \times 2^{34}$$

$$X = -1.1011 \times 2^{34} - 1.1111010 \times 2^{34} \text{ (unsigned addition)}$$

$$X = -1.1101001 \times 2^{34} = -1.1101001 \times 2^{35}$$

$$e + \text{bias} = 35 + 127 = 162 = 10100010$$

$$X = 1101 0001 0110 1001 0000 0000 0000 0000 = \text{D1690000}$$

$$\begin{array}{r} \begin{array}{cccccccc} b_7 & b_6 & b_5 & b_4 & b_3 & b_2 & b_1 & b_0 \end{array} \\ \begin{array}{cccccccc} 1 & 1 & 0 & 1 & 1 & 0 & 0 & + \\ 1 & 1 & 1 & 1 & 1 & 0 & 1 & \\ \hline 1 & 1 & 1 & 0 & 1 & 0 & 0 & 1 \end{array} \end{array}$$

- ✓ $X = 80C00000 \times \text{FAD00000}$:

80C00000: 1000 0000 1100 0000 0000 0000 0000 0000

$$e + \text{bias} = 00000001 = 1 \rightarrow e = 1 - 127 = -126$$

$$\text{Significand} = 1.1$$

$$80C00000 = -1.1 \times 2^{-126}$$

FAD00000: 1111 1010 1101 0000 0000 0000 0000 0000

$$e + \text{bias} = 11110101 = 245 \rightarrow e = 245 - 127 = 118$$

$$\text{Significand} = 1.101$$

$$\text{FAD00000} = -1.101 \times 2^{118}$$

$$X = (-1.1 \times 2^{-126}) \times (-1.101 \times 2^{118}) = 10.0111 \times 2^{-8} = 1.00111 \times 2^{-7}$$

$$e + \text{bias} = -7 + 127 = 120 = 01111000$$

$$X = 0011 1100 0001 1100 0000 0000 0000 0000 = \text{3C1C0000}$$

$$\begin{array}{r} \begin{array}{cccc} 1 & 1 & 0 & 1 \end{array} \times \\ \begin{array}{cccc} 1 & 1 & 1 & 1 \\ \hline 1 & 1 & 0 & 1 \\ \hline 1 & 0 & 0 & 1 & 1 & 1 \end{array} \end{array}$$

- ✓ $X = 7B380000 \div \text{C8A00000}$:

7B380000: 0111 1011 0011 1000 0000 0000 0000 0000

$$e + \text{bias} = 11110110 = 246 \rightarrow e = 246 - 127 = 119$$

$$\text{Significand} = 1.0111$$

$$7B380000 = 1.0111 \times 2^{119}$$

C8A00000: 1100 1000 1010 0000 0000 0000 0000 0000

$$e + \text{bias} = 10010001 = 145 \rightarrow e = 145 - 127 = 18$$

$$\text{Significand} = 1.01$$

$$\text{C8A00000} = -1.01 \times 2^{18}$$

$$X = -\frac{1.0111 \times 2^{119}}{1.01 \times 2^{18}} = -\frac{1.0111}{1.01} \times 2^{101}$$

$$\begin{array}{r}
 000010010 \\
 10100 \overline{) 101110000} \\
 \underline{10100} \\
 11000 \\
 \underline{10100} \\
 1000
 \end{array}$$

Alignment:

$$\frac{1.0111}{1.01} = \frac{1.0111}{1.0100} = \frac{10111}{10100}$$

Append $x = 4$ zeros: $\frac{101110000}{10100}$

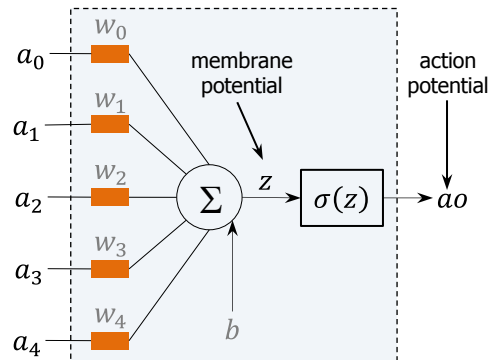
Integer division:
 $Q = 10010 \rightarrow Qf = 1.001$

Thus: $X = -1.001 \times 2^{101}$, $e + \text{bias} = 101 + 127 = 228 = 11100100$
 $X = 1111 \ 0010 \ 0001 \ 0000 \ 0000 \ 0000 \ 0000 \ 0000 = F2100000$

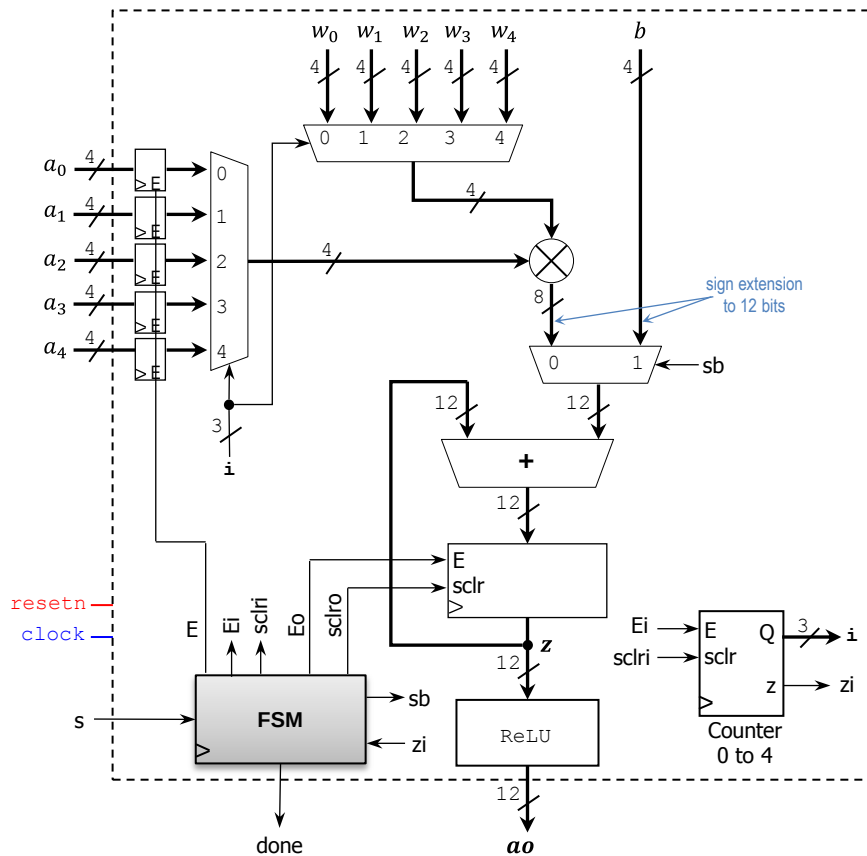
PROBLEM 4 (35 PTS)

- **Artificial neuron model.** The membrane potential z is a sum of products (input activations a_i by weights w_i) to which a bias term b is added. The action potential ao is modeled as a scalar function of z . The figure depicts a neuron with 5 inputs. The bias and the weights are constant values.

$$ao = \sigma \left(\sum_i a_i \times w_i + b \right)$$



- ✓ A popular and simple scalar function is the Rectified Linear Unit (ReLU):
 - $\sigma(z) = z$ if $z \geq 0$, otherwise $\sigma(z) = 0$.
- Digital System (FSM + Datapath): An iterative architecture for a 5-input neuron is depicted in next page. The circuit captures the input data (a_0, a_1, a_2, a_3, a_4) and then computes z using a multiply-and-accumulate approach (see iterative algorithm). The output ao is computed by applying the ReLU function to z .
 - ✓ All data is represented as signed integers:
 - Input activations (a_0, a_1, a_2, a_3, a_4), weights (w_0, w_1, w_2, w_3, w_4), bias (b): 4-bits wide.
 - Membrane potential (z) and action potential (ao): 12-bits wide (11 bits suffice, we select 12 for simplicity's sake).
 - ✓ Weights and biases: These are constant values (signed numbers represented as hexadecimals).
 - $w_0 = 0 \times 4$, $w_1 = 0 \times 1$, $w_2 = 0 \times 2$, $w_3 = 0 \times 8$, $w_4 = 0 \times 2$. $b = 0 \times 6$
 - ✓ Example:
 - If $a_0 = 0 \times 4$, $a_1 = 0 \times E$, $a_2 = 0 \times C$, $a_3 = 0 \times 5$, $a_4 = 0 \times A$. These values appear in the timing diagram (next page).
 - Then $z = 4 \times 4 + -2 \times 1 + -4 \times 2 + 5 \times (-8) + -6 \times 2 = -40 = 0 \times FD8$. Finally, $ao = 0 \times 000$
 - ✓ Counter 0 to 4: If $E=1$, $sclr=1$, then $Q \leftarrow 0$. If $E=1$, $sclr=0$, then $Q \leftarrow Q+1$. Also: $z=1$ if $Q = 4$, else $z=0$.
 - ✓ Register: If $E=1$, $sclr=1 \rightarrow$ Clear. If $E=1$, $sclr=0 \rightarrow$ Load data.
 - ✓ 4x4 Signed Multiplier: This is a combinational circuit, whose result is 8-bits wide (it should be sign-extended to 12 bits).
 - ✓ ReLU: Combinational Block that implements the ReLU operation. For example, if $z = 0 \times FD8$, then $ao = 0 \times 000$
- Sketch the Finite State Machine diagram (in ASM form) given the algorithm. (20 pts.)
 - ✓ The process begins when s is asserted, at this moment we capture a_0, a_1, a_2, a_3, a_4 on the input registers. Then z is updated until the counter reaches its maximum value (4). The signal done is asserted when the final result is computed.
- Complete the timing diagram (z and ao are in hexadecimal format). (15 pts.)

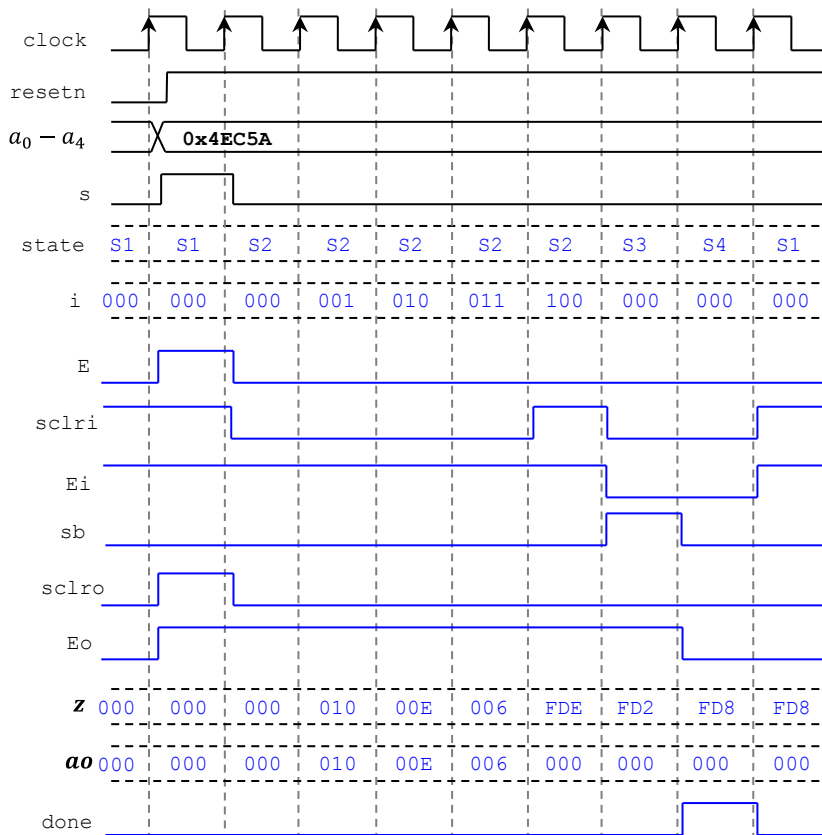


$$ao = \sigma_{ReLU} \left(\sum_{i=0}^4 a_i \times w_i + b \right)$$

ALGORITHM

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z ← 0
for i = 0 to 4
    z ← z + ai × wi
end
z ← z + b
ao ← z if z ≥ 0, else 0
    
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FSM:

